

Algebraic and combinatorial structures on Baxter permutations

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Abstract. We give a new construction of a Hopf subalgebra of the Hopf algebra of Free quasi-symmetric functions whose bases are indexed by objects belonging to the Baxter combinatorial family (*i.e.* Baxter permutations, pairs of twin binary trees, *etc.*). This construction relies on the definition of the Baxter monoid, analog of the plactic monoid and the sylvestre monoid, and on a Robinson-Schensted-like insertion algorithm. The algebraic properties of this Hopf algebra are studied. This Hopf algebra appeared for the first time in the work of Reading [Lattice congruences, fans and Hopf algebras, *Journal of Combinatorial Theory Series A*, 110:237–273, 2005].

Résumé. Nous proposons une nouvelle construction d’une sous-algèbre de Hopf de l’algèbre de Hopf des fonctions quasi-symétriques libres dont les bases sont indexées par les objets de la famille combinatoire de Baxter (*i.e.* permutations de Baxter, couples d’arbres binaires jumeaux, *etc.*). Cette construction repose sur la définition du monoïde de Baxter, analogue du monoïde plaxique et du monoïde sylvestre, et d’un algorithme d’insertion analogue à l’algorithme de Robinson-Schensted. Les propriétés algébriques de cette algèbre de Hopf sont étudiées. Cette algèbre de Hopf est apparue pour la première fois dans le travail de Reading [Lattice congruences, fans and Hopf algebras, *Journal of Combinatorial Theory Series A*, 110:237–273, 2005].

Keywords: Hopf algebras, Robinson-Schensted algorithm, quotient monoid, Baxter permutations

1 Introduction

In the recent years, many combinatorial Hopf algebras, whose bases are indexed by combinatorial objects, have been intensively studied. For example, the Malvenuto-Reutenauer Hopf algebra **FQSym** of Free quasi-symmetric functions [19, 7] has bases indexed by permutations. This Hopf algebra admits several Hopf subalgebras: The Hopf algebra of Free symmetric functions **FSym** [21, 7], whose bases are indexed by standard Young tableaux, the Hopf algebra **Bell** [23] whose bases are indexed by set partitions, the Loday-Ronco Hopf algebra **PBT** [18, 12] whose bases are indexed by planar binary trees and the Hopf algebra **Sym** of non-commutative symmetric functions [10] whose bases are indexed by integer compositions. An unifying approach to construct all these structures relies on a definition of a congruence on words leading to the definition of monoids on combinatorial objects. Indeed, **FSym** is directly obtained from the plactic monoid [15], **Bell** from the Bell monoid [23], **PBT** from the sylvestre monoid [11, 12], and **Sym** from the hypoplactic monoid [20]. The richness of these constructions relies on the fact that, in addition to construct Hopf algebras, the definition of such monoids often brings partial orders, combinatorial algorithms and Robinson-Schensted-like algorithms, of independent interest.

In this paper, we propose to enrich this collection of Hopf algebras by providing a construction of a Hopf algebra whose bases are indexed by objects belonging to the Baxter combinatorial family. This combinatorial family admits various representations as Baxter permutations [4], pairs of twin binary trees [8],

quadrangulations [1], plane bipolar orientations [5], *etc.* In [22], Reading defines first a Hopf algebra on Baxter permutations in the context of lattice congruences; Moreover, very recently, Law and Reading [16] have studied and detailed their construction of this Hopf algebra. However, even if both points of view lead to the same general theory, their paths are different and provide different ways of understanding this Hopf algebra, one centered, as in Law and Reading's work, on lattice theory, the other, as in our work, centered on combinatorics on words. Moreover, a large part of the results of each paper does not appear in the other.

We begin by recalling in Section 2 the preliminary notions used thereafter. In Section 3, we define the Baxter congruence. This congruence allows to define a quotient of the free monoid, the Baxter monoid, which has a number of properties required for the Hopf algebraic construction which follows. We show that the Baxter monoid is intimately linked to the sylvester monoid. Next, in Section 4, we develop a Robinson-Schensted-like insertion algorithm that allows to decide if two words are equivalent according to the Baxter congruence. Given a word, this algorithm computes a pair of twin binary trees. Section 5 is devoted to the study of some properties of the equivalence classes of permutations under the Baxter congruence. This leads to the definition of a lattice structure on pairs of twin binary trees. Finally, in Section 6, we define the Hopf algebra **Baxter** and study it. Using the order structure on pairs of twin binary trees, we provide multiplicative bases and show that **Baxter** is free as an algebra. Using the results of Foissy on bidendriform bialgebras [9], we show that **Baxter** is also self-dual and that the Lie algebra of its primitive elements is free.

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2 Preliminaries

2.1 Words

In the sequel, $A := \{a_1 < a_2 < \dots\}$ is a totally ordered infinite alphabet and A^* is the free monoid spanned by A . Let $u \in A^*$. For $S \subseteq A$, we denote by $u|_S$ the *restriction* of u on the alphabet S , that is the longest subword of u made of letters of S . The *evaluation* $\text{eval}(u)$ of the word u is the non-negative integer vector such that its i -th entry is the number of occurrences of the letter a_i in u . Let $\max(u)$ be the maximal letter of u . The *Schützenberger transformation* $\#$ is defined by $u^\# := \max(u)+1-u|_{\max(u)} \dots \max(u)+1-u_1$; For example, $(a_5 a_3 a_1 a_1 a_5 a_2)^\# = a_4 a_1 a_5 a_5 a_3 a_1$. Note that it is an involution if u has an occurrence of a_1 . Let $v \in A^*$ and $a, b \in A$. The *shuffle product* \sqcup is defined on $\mathbb{Z}\langle A \rangle$ recursively by $u \sqcup \epsilon := \epsilon \sqcup u := u$ and $au \sqcup bv := a(u \sqcup bv) + b(au \sqcup v)$.

2.2 Permutations

Denote by \mathfrak{S}_n the set of permutations of size n and $\mathfrak{S} := \cup_{n \geq 0} \mathfrak{S}_n$. We shall call (i, j) a *co-inversion* of $\sigma \in \mathfrak{S}$ if $i < j$ and $\sigma_i^{-1} > \sigma_j^{-1}$. Let us recall that the (*right*) *permutohedron order* is the partial order \leq_P defined on \mathfrak{S}_n where σ is covered by ν if $\sigma = uabv$ and $\nu = ubav$ where $a < b$. Let $\sigma, \nu \in \mathfrak{S}$. The permutation σ / ν is obtained by concatenating σ and the letters of ν incremented by $|\sigma|$; In the same way, the permutation $\sigma \setminus \nu$ is obtained by concatenating the letters of ν incremented by $|\sigma|$ and σ ; For example, $312 / 2314 = 3125647$ and $312 \setminus 2314 = 5647312$. The permutation σ is *connected* if $\sigma = \nu / \pi$ implies $\nu = \sigma$ or $\pi = \sigma$. The *shifted shuffle product* \sqcup of two permutations is defined by $\sigma \sqcup \nu :=$

$\sigma \sqcup (\nu_1 + |\sigma| \dots \nu_{|\nu|} + |\sigma|)$; For example, $\mathbf{12} \sqcup \mathbf{21} = \mathbf{12} \sqcup \mathbf{43} = \mathbf{1243} + \mathbf{1423} + \mathbf{1432} + \mathbf{4123} + \mathbf{4132} + \mathbf{4312}$. The *standardized word* $\text{std}(u)$ of $u \in A^*$ is the unique permutation σ satisfying $\sigma_i < \sigma_j$ iff $u_i \leq u_j$ for all $1 \leq i < j \leq |u|$; For example, $\text{std}(a_3a_1a_4a_2a_5a_7a_4a_2a_3) = 416289735$.

2.3 Binary trees

Denote by \mathcal{BT}_n the set of binary trees with n internal nodes and $\mathcal{BT} := \cup_{n \geq 0} \mathcal{BT}_n$. We use in the sequel the standard terminology (i.e., *child*, *ancestor*, ...) about binary trees [2]. The only element of \mathcal{BT}_0 is the *leaf* or *empty tree*, denoted by \perp . Let us recall that the *Tamari order* [14] is the partial order \leq_T defined on \mathcal{BT}_n where $T_0 \in \mathcal{BT}_n$ is covered by $T_1 \in \mathcal{BT}_n$ if it is possible to transform T_0 into T_1 by performing a right rotation (see Figure 1).

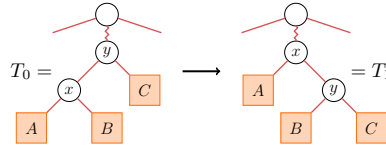


Figure 1: The right rotation of root y .

Let $T_0, T_1 \in \mathcal{BT}$. The binary tree $T_0 \diagup T_1$ is obtained by grafting T_0 from its root on the leftmost leaf of T_1 ; In the same way, the binary tree $T_0 \diagdown T_1$ is obtained by grafting T_1 from its root on the rightmost leaf of T_0 . The *canopy* (see [18] and [26]) $\text{cnp}(T)$ of $T \in \mathcal{BT}$ is the word on the alphabet $\{0, 1\}$ obtained by browsing the leaves of T from left to right except the first and the last one, writing 0 if the considered leaf is oriented to the right, 1 otherwise (see Figure 2). Note that the orientation of the leaves in a binary tree is determined only by its nodes so that we can omit to draw the leaves in our next graphical representations.

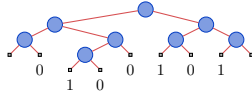


Figure 2: The canopy of this binary tree is 0100101.

An A -labeled binary tree T is a *left* (resp. *right*) *binary search tree* if for any node x labeled by b , each label a of a node in the left subtree of x and each label c of a node in the right subtree of x , the inequality $a < b \leq c$ (resp. $a \leq b < c$) holds. A binary tree $T \in \mathcal{BT}_n$ is a *decreasing binary tree* if it is bijectively labeled on $\{1, \dots, n\}$ and, for all node y of T , if x is a child of y , then the label of x is smaller than the label of y . The *shape* of a labeled binary tree is the unlabeled binary tree obtained by forgetting its labels.

2.4 Baxter permutations and pairs of twin binary trees

A permutation σ is a *Baxter permutation* if for any subword $u = u_1u_2u_3u_4$ of σ such that the letters u_2 and u_3 are adjacent in σ , $\text{std}(u) \notin \{2413, 3142\}$. In other words, σ is a Baxter permutation if it avoids the *generalized permutation patterns* $2 - 41 - 3$ and $3 - 14 - 2$ (see [3] for an introduction on generalized permutation patterns). For example, 42173856 is not a Baxter permutation; On the other hand 436975128 is a Baxter permutation. Let us denote by \mathfrak{S}_n^B the set of Baxter permutations of size n and $\mathfrak{S}^B := \cup_{n \geq 0} \mathfrak{S}_n^B$.

A *pair of twin binary trees* (T_L, T_R) is made of two binary trees $T_L, T_R \in \mathcal{BT}_n$ such that the canopies of T_L and T_R are complementary, that is $\text{cnp}(T_L)_i \neq \text{cnp}(T_R)_i$ for all $1 \leq i \leq n - 1$. Denote by \mathcal{TB}_n

the set of pairs of twin binary trees where each binary tree has n nodes and $\mathcal{TB}\mathcal{T} := \cup_{n \geq 0} \mathcal{TB}\mathcal{T}_n$. In [8], Dulucq and Guibert have highlighted a bijection between Baxter permutations and pairs of twin binary trees. In the sequel, we shall make use of a very similar bijection.

3 The Baxter monoid

3.1 Definition and first properties

Recall that an equivalence relation \equiv defined on A^* is a *congruence* if for all $u, u', v, v' \in A^*$, $u \equiv u'$ and $v \equiv v'$ imply $u.v \equiv u'.v'$.

Definition 3.1 *The Baxter monoid is the quotient of the free monoid A^* by the congruence \equiv_B that is the transitive closure of the adjacency relations \leftrightsquigarrow_B and \rightrightarrows_B defined for $u, v \in A^*$ and $a, b, c, d \in A$ by:*

$$cuadvb \leftrightsquigarrow_B cudavb \quad \text{where} \quad a \leq b < c \leq d, \quad (1)$$

$$budavc \rightrightarrows_B buadv c \quad \text{where} \quad a < b \leq c < d. \quad (2)$$

For $u \in A^*$, denote by \hat{u} the \equiv_B -equivalence class of u ; For example, the \equiv_B -equivalence class of 5273641 is $\{5237641, 5273641, 5276341, 5723641, 5726341, 5762341\}$.

An equivalence relation \equiv defined on A^* is *compatible with the restriction of alphabet intervals* if for all interval I of A and for all $u, v \in A^*$, $u \equiv v$ implies $u|_I \equiv v|_I$.

Proposition 3.2 *The Baxter monoid is compatible with the restriction of alphabet intervals.*

Proof: We only have to check the property on adjacency relations. □

An equivalence relation \equiv defined on A^* is *compatible with the destandardization process* if for all $u, v \in A^*$, $u \equiv v$ iff $\text{std}(u) \equiv \text{std}(v)$ and $\text{eval}(u) = \text{eval}(v)$.

Proposition 3.3 *The Baxter monoid is compatible with the destandardization process.*

An equivalence relation \equiv defined on A^* is *compatible with the Schützenberger involution* if for all $u, v \in A^*$, $u \equiv v$ implies $u^\# \equiv v^\#$.

Proposition 3.4 *The Baxter monoid is compatible with the Schützenberger involution.*

3.2 Connection with the sylvester monoid

The *sylvester monoid* [11, 12] is the quotient of the free monoid A^* by the congruence \equiv_S that is the transitive closure of the adjacency relation \leftrightsquigarrow_S defined for $u \in A^*$ and $a, b, c \in A$ by:

$$acub \leftrightsquigarrow_S caub \quad \text{where} \quad a \leq b < c. \quad (3)$$

In the same way, let us define the *#-sylvester monoid* by the congruence $\equiv_{S^\#}$ that is the transitive closure of the adjacency relation $\leftrightsquigarrow_{S^\#}$ defined for $u \in A^*$ and $a, b, c \in A$ by:

$$buac \leftrightsquigarrow_{S^\#} buca \quad \text{where} \quad a < b \leq c. \quad (4)$$

Note that this adjacency relation is defined by taking the images by the Schützenberger involution of the sylvester adjacency relation. Indeed, for all $u, v \in A^*$, $u \equiv_{S^\#} v$ iff $u^\# \equiv_S v^\#$. The Baxter monoid and the sylvester monoid are related in the following way:

Proposition 3.5 *Let $u, v \in A^*$. Then, $u \equiv_B v$ iff $u \equiv_S v$ and $u \equiv_{S^\#} v$.*

Proposition 3.5 shows that the \equiv_B -equivalence classes are the intersection of \equiv_S -equivalence classes and $\equiv_{S^\#}$ -equivalence classes.

4 A Robinson-Schensted-like algorithm

We shall describe here an insertion algorithm $u \mapsto (\mathbb{P}(u), \mathbb{Q}(u))$, such that, given a word $u \in A^*$, it computes its \mathbb{P} -symbol, that is a pair of A -labeled twin binary trees (T_L, T_R) where T_L (resp. T_R) is a left (resp. right) binary search tree, and its \mathbb{Q} -symbol, a decreasing binary tree.

4.1 Definition of the insertion algorithm

Let T be an A -labeled right binary search tree and b a letter of A . The *lower restricted binary tree* of T compared to b , namely $T_{\leq b}$, is the right binary search tree uniquely made of the nodes x of T labeled by a letter a satisfying $a \leq b$ and such that for all nodes x and y of $T_{\leq b}$, if x is ancestor of y in $T_{\leq b}$, then x is ancestor of y in T . In the same way, we define the *higher restricted binary tree* of T compared to b , namely $T_{>b}$ (see Figure 3).

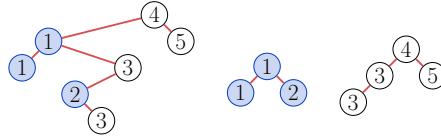


Figure 3: A right binary search tree T , $T_{\leq 2}$ and $T_{>2}$.

Let T be an A -labeled right binary search tree and a a letter of A . The *root insertion* of a into T consists in modifying T so that the root of T is a new node labeled by a , its left subtree is $T_{\leq a}$ and its right subtree is $T_{>a}$.

Let T be an A -labeled left (resp. right) binary search tree and a a letter of A . The *leaf insertion* of a into T is recursively defined by: If $T = \perp$, the result is the one-node binary tree labeled by a ; Else, if the label b of the root of T satisfies $a < b$ (resp. $a \leq b$), make a leaf insertion of a into the left subtree of T , else, make a leaf insertion of a into the right subtree of T .

Given a pair of A -labeled twin binary trees (T_L, T_R) where T_L (resp. T_R) is a left (resp. right) binary search tree, the *insertion* of the letter a of A into (T_L, T_R) consists in making a leaf insertion of a into T_L and a root insertion of a into T_R .

The \mathbb{P} -symbol (T_L, T_R) of a word $u \in A^*$ is computed by iteratively inserting the letters of u , from left to right, into the pair of twin binary trees (\perp, \perp) . The \mathbb{Q} -symbol of u is the decreasing binary tree labeled on $\{1, \dots, |u|\}$, built by recording the dates of creation of each node of T_R (see Figure 4).

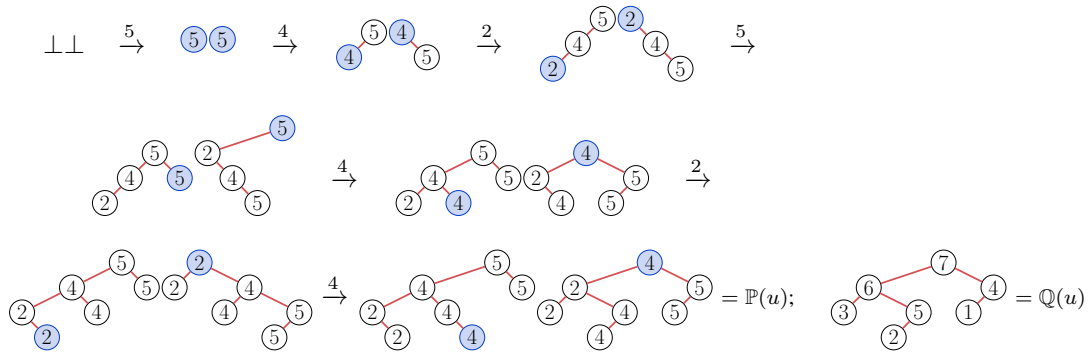


Figure 4: Steps of computation of the \mathbb{P} -symbol and the \mathbb{Q} -symbol of $u := 5425424$.

4.2 Validity of the insertion algorithm

Lemma 4.1 *Let $u \in A^*$. Let T be the right binary search tree obtained by root insertions of the letters of u , from left to right. Let T' be the right binary search tree obtained by leaf insertions of the letters of u , from right to left. Then, $T = T'$.*

Lemma 4.2 *Let $\sigma \in \mathfrak{S}$ and $T \in \mathcal{BT}_{|\sigma|}$ be the binary search tree obtained by leaf insertions of the letters of σ , from left to right. Then, for $1 \leq i \leq |\sigma| - 1$, the $i + 1$ -st leaf of T is right-oriented iff $(i, i + 1)$ is a co-inversion of σ .*

If (T_L, T_R) is a pair of labeled twin binary trees, define its *shape*, that is the pair of unlabeled twin binary trees (T'_L, T'_R) where T'_L (resp. T'_R) is the shape of T_L (resp. T_R).

Proposition 4.3 *For all word $u \in A^*$, the shape of the \mathbb{P} -symbol of u is a pair of twin binary trees.*

Proposition 4.4 *Let $u, v \in A^*$. Then, $u \equiv_B v$ iff $\mathbb{P}(u) = \mathbb{P}(v)$.*

In particular, we have $\mathbb{P}(\sigma) = \mathbb{P}(\nu)$ iff the permutations σ and ν are \equiv_B -equivalent. Moreover, each \equiv_B -equivalence class of permutations can be encoded by a pair of unlabeled twin binary trees because there is one unique way to bijectively label a binary tree with n nodes on $\{1, \dots, n\}$ such that it is a binary search tree.

Remark 4.5 *Let $u, v \in A^*$ and $(T_L, T_R) := \mathbb{P}(u)$. We have $u \equiv_B v$ iff the following two assertions are satisfied:*

- (i) *v is a linear extension of T_L seen as a poset in which the smallest element is its root;*
- (ii) *v is a linear extension of T_R seen as a poset in which minimal elements are the nodes with no descents.*

5 The Baxter lattice

5.1 Some properties of the \equiv_B -equivalence classes of permutations

Theorem 5.1 *For all $n \geq 0$, each equivalence class of $\mathfrak{S}_n / \equiv_B$ contains exactly one Baxter permutation.*

Proposition 5.2 *For all $n \geq 0$, each equivalence class of $\mathfrak{S}_n / \equiv_B$ is an interval of the permutohedron.*

For all permutation σ , let us define $\sigma \uparrow$ (resp. $\sigma \downarrow$) the maximal (resp. minimal) permutation of the \equiv_B -equivalence class of σ for the permutohedron order.

Proposition 5.3 *Let $\sigma, \nu \in \mathfrak{S}_n$ such that $\sigma \leq_P \nu$. Then, $\sigma \uparrow \leq_P \nu \uparrow$ and $\sigma \downarrow \leq_P \nu \downarrow$.*

5.2 A lattice structure on the set of pairs of twin binary trees

Definition 5.4 *For all $n \geq 0$, define the order relation \leq_B on the set $\mathcal{TB}\mathcal{T}_n$ setting $J_0 \leq_B J_1$, where $J_0, J_1 \in \mathcal{TB}\mathcal{T}_n$, if there exists $\sigma_0, \sigma_1 \in \mathfrak{S}_n$ such that $\mathbb{P}(\sigma_0) = J_0$, $\mathbb{P}(\sigma_1) = J_1$ and $\sigma_0 \leq_P \sigma_1$.*

Propositions 5.2 and 5.3 ensure that this order is well-defined, and in particular that the relation \leq_B is transitive and antisymmetric.

The pair of twin binary trees (T_L, T_R) is covered by $(T'_L, T'_R) \in \mathcal{TB}\mathcal{T}$ if one of the three following conditions is satisfied:

1. $T'_R = T_R$ and T'_L is obtained from T_L by performing a left rotation into T_L such that $\text{cnp}(T_L) = \text{cnp}(T'_L)$;

2. $T'_L = T_L$ and T'_R is obtained from T_R by performing a right rotation into T_R such that $\text{cnp}(T_R) = \text{cnp}(T'_R)$;
3. T'_L (resp. T'_R) is obtained by performing a left (resp. right) rotation into T_L (resp. T_R) such that $\text{cnp}(T_L) \neq \text{cnp}(T'_L)$ (resp. $\text{cnp}(T_R) \neq \text{cnp}(T'_R)$).

Moreover, it is possible to compare two pairs of twin binary trees $J_0 := (T_L^0, T_R^0)$ and $J_1 := (T_L^1, T_R^1)$ very easily by computing the *Tamari vector* (see [14]) of each binary tree. Indeed, we have $J_0 \leq_B J_1$ iff the Tamari vector of T_L^0 (resp. T_R^0) is greater (resp. smaller) component by component than the Tamari vector of T_L^1 (resp. T_R^1).

Propositions 5.2 and 5.3 implies that \equiv_B is also a lattice congruence [6, 22]. Thus, since the permutohedron is a lattice,

Proposition 5.5 *For all $n \geq 0$, the poset (\mathcal{BT}_n, \leq_B) is a lattice.*

6 The Baxter Hopf Algebra

In the sequel, all the algebraic structures have a field of characteristic zero \mathbb{K} as ground field.

6.1 The Hopf algebra \mathbf{FQSym}

Recall that the family $\{\mathbf{F}_\sigma\}_{\sigma \in \mathfrak{S}}$ form the *fundamental basis* of \mathbf{FQSym} [7]. Its product and its coproduct are defined by:

$$\mathbf{F}_\sigma \cdot \mathbf{F}_\nu := \sum_{\pi \in \sigma \sqcup \nu} \mathbf{F}_\pi, \quad \Delta(\mathbf{F}_\sigma) := \sum_{0 \leq i \leq |\sigma|} \mathbf{F}_{\text{std}(\sigma_1 \dots \sigma_i)} \otimes \mathbf{F}_{\text{std}(\sigma_{i+1} \dots \sigma_{|\sigma|})}. \quad (5)$$

The following theorem due to Hivert and Nzeutchap [13] shows that an equivalence relation on A^* satisfying some properties can be used to define Hopf subalgebras of \mathbf{FQSym} :

Theorem 6.1 *Let \equiv be an equivalence relation defined on A^* . If \equiv is a congruence, compatible with the restriction of alphabet intervals and compatible with the destandardization process, then, the family $\{\mathbf{P}_{\hat{\sigma}}\}_{\hat{\sigma} \in \mathfrak{S}/\equiv}$ defined by:*

$$\mathbf{P}_{\hat{\sigma}} := \sum_{\sigma \in \hat{\sigma}} \mathbf{F}_\sigma \quad (6)$$

spans a Hopf subalgebra of \mathbf{FQSym} .

6.2 The Hopf algebra \mathbf{Baxter}

By definition, \equiv_B is a congruence, and, by Proposition 3.2 and 3.3, \equiv_B checks the conditions of Theorem 6.1. Moreover, by Proposition 4.4, the \equiv_B -equivalence classes of permutations can be encoded by pairs of unlabeled twin binary trees. Hence, we have the following theorem:

Theorem 6.2 *The family $\{\mathbf{P}_J\}_{J \in \mathcal{BT}}$ defined by:*

$$\mathbf{P}_J := \sum_{\substack{\sigma \in \mathfrak{S} \\ \mathbb{P}(\sigma) = J}} \mathbf{F}_\sigma \quad (7)$$

spans a Hopf subalgebra of \mathbf{FQSym} , namely the Hopf algebra \mathbf{Baxter} .

The Hilbert series of **Baxter** is $B(z) := 1 + z + 2z^2 + 6z^3 + 22z^4 + 92z^5 + 422z^6 + 2074z^7 + 10754z^8 + 58202z^9 + 326240z^{10} + 1882960z^{11} + \dots$, the generating series of Baxter permutations (sequence **A001181** of [24]).

One has for example,

$$\mathbf{P}_{\text{[diagram]}} = \mathbf{F}_{12}, \quad \mathbf{P}_{\text{[diagram]}} = \mathbf{F}_{2143} + \mathbf{F}_{2413}, \quad \mathbf{P}_{\text{[diagram]}} = \mathbf{F}_{542163} + \mathbf{F}_{542613} + \mathbf{F}_{546213}. \quad (8)$$

By Theorem 6.1, the product of **Baxter** is well-defined. We deduce it from the product of **FQSym** and obtain

$$\mathbf{P}_{J_0} \cdot \mathbf{P}_{J_1} = \sum_{\substack{\mathbb{P}(\sigma)=J_0, \mathbb{P}(\nu)=J_1 \\ \pi \in \sigma \sqcup \nu \cap \mathfrak{S}^{\mathbf{B}}}} \mathbf{P}_{\mathbb{P}(\pi)}. \quad (9)$$

For example,

$$\begin{aligned} \mathbf{P}_{\text{[diagram]}} \cdot \mathbf{P}_{\text{[diagram]}} = & \mathbf{P}_{\text{[diagram]}} + \mathbf{P}_{\text{[diagram]}} + \mathbf{P}_{\text{[diagram]}} \\ & + \mathbf{P}_{\text{[diagram]}} + \mathbf{P}_{\text{[diagram]}} + \mathbf{P}_{\text{[diagram]}}. \end{aligned} \quad (10)$$

In the same way, we deduce the coproduct of **Baxter** from the coproduct of **FQSym** and obtain

$$\Delta(\mathbf{P}_J) = \sum_{\substack{\mathbb{P}(\pi)=J \\ \pi=uv \\ \sigma:=\text{std}(u), \nu:=\text{std}(v) \in \mathfrak{S}^{\mathbf{B}}}} \mathbf{P}_{\mathbb{P}(\sigma)} \otimes \mathbf{P}_{\mathbb{P}(\nu)}. \quad (11)$$

For example,

$$\begin{aligned} \Delta \mathbf{P}_{\text{[diagram]}} = & 1 \otimes \mathbf{P}_{\text{[diagram]}} + \mathbf{P}_{\text{[diagram]}} \otimes \mathbf{P}_{\text{[diagram]}} + \mathbf{P}_{\text{[diagram]}} \otimes \mathbf{P}_{\text{[diagram]}} + \mathbf{P}_{\text{[diagram]}} \otimes \mathbf{P}_{\text{[diagram]}} \\ & + \mathbf{P}_{\text{[diagram]}} \otimes \mathbf{P}_{\text{[diagram]}} + \mathbf{P}_{\text{[diagram]}} \otimes \mathbf{P}_{\text{[diagram]}} + \mathbf{P}_{\text{[diagram]}} \otimes \mathbf{P}_{\text{[diagram]}} + \mathbf{P}_{\text{[diagram]}} \otimes 1. \end{aligned} \quad (12)$$

Remark 6.3 It is well-known that the Hopf algebra **PBT** [18, 12] is a Hopf subalgebra of **FQSym**. Besides, we have the following sequence of injective Hopf maps:

$$\mathbf{PBT} \xhookrightarrow{\rho} \mathbf{Baxter} \hookrightarrow \mathbf{FQSym}. \quad (13)$$

Indeed, by Proposition 3.5, every $\equiv_{\mathbf{S}}$ -equivalence class is an union of some $\equiv_{\mathbf{B}}$ -equivalence classes. Denoting by $\{\mathbf{P}_T\}_{T \in \mathbf{BT}}$ the basis of **PBT** defined in accordance with (6) by the sylvester equivalence relation $\equiv_{\mathbf{S}}$, we have

$$\rho(\mathbf{P}_T) = \sum_{\substack{T' \in \mathbf{BT} \\ J := (T', T) \in \mathbf{TBT}}} \mathbf{P}_J. \quad (14)$$

For example,

$$\rho\left(\mathbf{P}_{\text{[diagram]}}\right) = \mathbf{P}_{\text{[diagram]}} + \mathbf{P}_{\text{[diagram]}} + \mathbf{P}_{\text{[diagram]}}. \quad (15)$$

6.3 Multiplicative bases

Define the *elementary* family $\{\mathbf{E}_J\}_{J \in \mathcal{TB}\mathcal{T}}$ and the *homogeneous* family $\{\mathbf{H}_J\}_{J \in \mathcal{TB}\mathcal{T}}$ respectively by:

$$\mathbf{E}_J := \sum_{J \leq_B J'} \mathbf{P}_{J'} \quad \text{and} \quad \mathbf{H}_J := \sum_{J' \leq_B J} \mathbf{P}_{J'}. \quad (16)$$

These families are bases of **Baxter** since they are defined by triangularity.

Let $J_0 := (T_L^0, T_R^0)$ and $J_1 := (T_L^1, T_R^1)$ be two pairs of twin binary trees. Let us define the pair of twin binary trees J_0 / J_1 by $J_0 / J_1 := (T_L^0 \setminus T_L^1, T_R^0 / T_R^1)$. In the same way, the pair of twin binary trees $J_0 \setminus J_1$ is defined by $J_0 \setminus J_1 := (T_L^0 / T_L^1, T_R^0 \setminus T_R^1)$.

Using the multiplicative bases of **FQSym**, we establish the following proposition:

Proposition 6.4 *For all $J_0, J_1 \in \mathcal{TB}\mathcal{T}$, we have*

$$\mathbf{E}_{J_0} \cdot \mathbf{E}_{J_1} = \mathbf{E}_{J_0 / J_1} \quad \text{and} \quad \mathbf{H}_{J_0} \cdot \mathbf{H}_{J_1} = \mathbf{H}_{J_0 \setminus J_1}. \quad (17)$$

Lemma 6.5 *Let C be an equivalence class of $\mathfrak{S}_n / \equiv_B$. The Baxter permutation belonging to C is connected iff all the permutations of C are connected.*

Let us say that a pair of twin binary trees J is *connected* if the unique Baxter permutation σ satisfying $\mathbb{P}(\sigma) = J$ is connected.

Proposition 6.6 *The Hopf algebra **Baxter** is free on the elements \mathbf{E}_J where J is a connected pair of twin binary trees.*

The generating series $B_C(z)$ of connected Baxter permutations is $B_C(z) = 1 - B(z)^{-1}$. First dimensions of algebraic generators of **Baxter** are 1, 1, 1, 3, 11, 47, 221, 1113, 5903, 32607, 186143, 1092015.

6.4 Bidendriform bialgebra structure

A Hopf algebra (H, \cdot, Δ) can be fitted into a bidendriform bialgebra structure [9] if (H^+, \prec, \succ) is a dendriform algebra [17] and $(H^+, \Delta_\prec, \Delta_\succ)$ a codendriform coalgebra, where H^+ is the augmentation ideal of H . The operators $\prec, \succ, \Delta_\prec$ and Δ_\succ have to fulfil some compatibility relations. In particular, for all $x, y \in H^+$, the product \cdot of H is retrieved by $x \cdot y = x \prec y + x \succ y$ and the coproduct Δ of H is retrieved by $\Delta(x) = 1 \otimes x + \Delta_\prec(x) + \Delta_\succ(x) + x \otimes 1$.

The Hopf algebra **FQSym** admits a bidendriform bialgebra structure [9]. Indeed, for all $\sigma, \nu \in \mathfrak{S}$ set

$$\mathbf{F}_\sigma \prec \mathbf{F}_\nu := \sum_{\substack{\pi \in \sigma \sqcup \nu \\ \pi|_{\pi|} = \sigma|_{\sigma|}}} \mathbf{F}_\pi, \quad \mathbf{F}_\sigma \succ \mathbf{F}_\nu := \sum_{\substack{\pi \in \sigma \sqcup \nu \\ \pi|_{\pi|} = \nu|_{\nu|} + |\sigma|}} \mathbf{F}_\pi, \quad (18)$$

$$\Delta_\prec(\mathbf{F}_\sigma) := \sum_{\substack{\sigma|_{\sigma|}^{-1} \leq i \leq |\sigma| - 1}} \mathbf{F}_{\text{std}(\sigma_1 \dots \sigma_i)} \otimes \mathbf{F}_{\text{std}(\sigma_{i+1} \dots \sigma_{|\sigma|})}, \quad (19)$$

$$\Delta_\succ(\mathbf{F}_\sigma) := \sum_{1 \leq i \leq \sigma|_{\sigma|}^{-1} - 1} \mathbf{F}_{\text{std}(\sigma_1 \dots \sigma_i)} \otimes \mathbf{F}_{\text{std}(\sigma_{i+1} \dots \sigma_{|\sigma|})}. \quad (20)$$

Proposition 6.7 *If \equiv is an equivalence relation defined on A^* satisfying the conditions of Theorem 6.1 and additionally, for all $u, v \in A^*$, the relation $u \equiv v$ implies $u|_u = v|_v$, then, the family defined in (6) spans a bidendriform sub-bialgebra of **FQSym**, and is free as an algebra, cofree as a coalgebra, self-dual, and the Lie algebra of its primitive elements is free.*

The equivalence relation \equiv_B satisfies the premises of Proposition 6.7 so that **Baxter** is free as an algebra, cofree as a coalgebra, self-dual, and the Lie algebra of its primitive elements is free.

6.6.2 Totally primitive elements

An element x of a bidendriform bialgebra is *totally primitive* if $\Delta_{\prec}(x) = 0 = \Delta_{\succ}(x)$.

Following [9], the generating series $B_T(z)$ of the totally primitive elements of **Baxter** is $B_T(z) = \frac{B(z)-1}{B(z)^2}$. First dimensions of totally primitive elements of **Baxter** are 0, 1, 0, 1, 4, 19, 96, 511, 2832, 16215, 95374, 573837. Here follows a basis of the totally primitive elements of **Baxter** of order 1, 3 and 4:

$$t_{1,1} = \mathbf{P} \bullet \bullet, \quad (28)$$

$$t_{3,1} = \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \\ \bullet \bullet \bullet \end{array} - \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \\ \bullet \bullet \bullet \end{array}, \quad (29)$$

$$t_{4,1} = \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \end{array} + \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \end{array} + \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \end{array} + \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \end{array} - \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \end{array} - \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \end{array} - \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \end{array}, \quad (30)$$

$$t_{4,2} = \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \end{array} - \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \end{array}, \quad (31)$$

$$t_{4,3} = \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \end{array} - \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \end{array}, \quad (32)$$

$$t_{4,4} = \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \end{array} - \mathbf{P} \begin{array}{c} \bullet \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \end{array}. \quad (33)$$

Baxter is free as dendriform algebra on its totally primitive elements.

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